

2. (18 pts) ON FA AND REGULAR LANGUAGES
Say whether each of the following languages is regular or not regular? Prove your answers.

(a) $L_1 = \{uvu^R \mid u \in \{a,b\}^+, \text{ and } v \in \{a,b\}\}$.

(b) $L_2 = \{wv^R \mid w \in \{a,b\}^+, \text{ and } v \in \{a,b\}^+\}$.

3. (18 pts) ON PDA AND CONTEXT-FREE LANGUAGES

Let $L_3 = \{xx^Ryy^R \mid x,y \in \{a,b\}^+, |y| \text{ is odd, and } bba \text{ is a substring of } x\}$.

- (a) Construct a context-free grammar that generates the language L_3 .
- (b) Construct a pushdown automata that accepts L_3 .

SOLUTION: (a)

(b) The PDA $M = (K, \Sigma, \Gamma, \Delta, s, F)$ is defined below:

it is true fill a T otherwise a F in the bracket before the statement.

- (a) () For any regular languages $L_1 \subseteq L_2 \subseteq \dots \subseteq L_n \subseteq \dots$, then the union $\bigcup_{n=1}^{\infty} L_n$ is also regular.
- (b) () If A is regular and both of B and $A \cap B$ are non-regular, then $A \cup B$ is non-regular.
- (c) () Let $D_{DFA} = \{ \langle M \rangle \mid M \text{ is a DFA, } \langle M \rangle \notin L(M) \}$, then D_{DFA} is not regular, but recursive, where " M " is the encoding of DFA M , just as Turing Machine.
- (d) () The language $\{a^i b^j c^k \mid i, j, k \in \mathbb{N}, \text{ and } k \neq i + j\}$ is context-free.
- (e) () A is recursively enumerable and B is regular, then $A \cap B$ is recursive.
- (f) () Let A and B be two recursively enumerable language. If both $\overline{A \cup B}$ and $\overline{A \cap B}$ are also recursively enumerable, then A and B are recursive.
- (g) () Let $A = \{ \langle M \rangle \mid \text{TM } M \text{ halts on at least 2020 strings} \}$. Suppose $A \leq_r \overline{B}$, then B is not recursively enumerable.
- (h) () Let A, B, C be arbitrary languages. If $A \leq C, B \leq C$ and C is recursively enumerable, then $A \cap B$ is recursively enumerable.
- (i) () Let A be a language, if there is a Turing machine M halts on x for every $x \in A$, then A is recursive
- (j) () If A is a Turing-enumerable language, then either A is recursive, or \overline{A} is not Turing-enumerable.
- (k) () There are some languages in NP are not recursive.
- (l) () There are some languages that cannot be semi-decided by any Turing machine.

	(q, a, b)	(q, γ)
$K =$ _____		
$\Sigma = \{a, b\}$		
$\Gamma =$ _____		
$s =$ _____		
$F =$ _____		

4. (12 pts) Construct a Turing machine that decides the following language:

$$L_4 = \{ucvuw^R \mid u, v, w \in \{a, b\}^*\}$$

When describing the Turing machines above, you can use the elementary Turing machines described in textbook. Always assume that the Turing machines start computation from the configuration $\triangleright _ _ x$ where $x \in \{a, b, c\}^*$ is the input string.

5. (10 pts) ON PRIMITIVE RECURSIVE FUNCTION

Show the following function $\varphi : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ by

$$\varphi(x, y) = \begin{cases} x/2 + 2y, & \text{if } x \text{ is even, and } y \text{ is a composite number} \\ |x - y|, & \text{otherwise} \end{cases}$$

is primitive recursive.